

# Basic Discrete Mathematics

## Review 2

**Meng Zhang**

ZJU-UIUC Institute  
Zhejiang University  
Email: mengzhang@intl.zju.edu.cn



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# Lecture Schedule

- |   |                                |   |           |
|---|--------------------------------|---|-----------|
| 4 | Number Theory and Cryptography | 7 | Counting  |
| 5 | Mathematical Induction         |   |           |
| 6 | Recursion                      | 8 | Relations |



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- 4 **Number Theory and Cryptography**
- 5 Mathematical Induction
- 6 Recursion
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# GCD as Linear Combinations

**Bezout's Theorem:** If  $a$  and  $b$  are positive integers, then there exist integers  $s$  and  $t$  such that

$$\gcd(a, b) = sa + tb.$$

This equation is called Bezout's identity.

We can use **extended Euclidean algorithm** to find Bezout's identity.

**Lemma:** If  $a, b, c$  are positive integers such that  $\gcd(a, b) = 1$  and  $a|bc$ , then  $a|c$ .

**Lemma:** If  $p$  is prime and  $p|a_1a_2\dots a_n$ , then  $p|a_i$  for some  $i$ .



# Linear Congruences

A congruence of the form  $ax \equiv b \pmod{m}$ , where  $m$  is a positive integer,  $a$  and  $b$  are integers, and  $x$  is a variable, is called a **linear congruence**.

The solutions to a linear congruence  $ax \equiv b \pmod{m}$  are **all integers  $x$**  that satisfy the congruence.

**Modular Inverse:** An integer  $\bar{a}$  such that  $\bar{a}a \equiv 1 \pmod{m}$  is said to be an **inverse** of  $a$  modulo  $m$ .

Solve the congruence  $ax \equiv b \pmod{m}$  by **multiplying both sides by  $\bar{a}$** .

$$x \equiv \bar{a}b \pmod{m}.$$



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# Modular Inverse

**Modular Inverse:** An integer  $\bar{a}$  such that  $\bar{a}a \equiv 1 \pmod{m}$  is said to be an **inverse** of  $a$  modulo  $m$ .

When does inverse exist?

**Theorem:** If  $a$  and  $m$  are **relatively prime integers** and  $m > 1$ , then an inverse of  $a$  modulo  $m$  **exists**. The inverse is **unique** modulo  $m$ . That is,

- there is a unique positive integer  $\bar{a}$  less than  $m$  that is an inverse of  $a$  modulo  $m$  and
- every other inverse of  $a$  modulo  $m$  is congruent to  $\bar{a}$  modulo  $m$ .

If we obtain an arbitrary inverse of  $a$  modulo  $m$ , how to obtain the inverse that is less than  $m$ ?



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# Modular Inverse

## How to find inverses?

Using **extended Euclidean algorithm**:

**Example:** Find an inverse of 101 modulo 4620. That is, find  $\bar{a}$  such that  $\bar{a} \cdot 101 \equiv 1 \pmod{4620}$ .

With extended Euclidean algorithm, we obtain  $\gcd(a, b) = sa + tb$ , i.e.,  $1 = -35 \cdot 4620 + 1601 \cdot 101$ . It tells us that  $-35$  and  $1601$  are Bezout coefficients of 4620 and 101. We have

$$1 \pmod{4620} = 1601 \cdot 101 \pmod{4620}.$$

Thus, 1601 is an inverse of 101 modulo 4620.



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# The Chinese Remainder Theorem

**Theorem** (The Chinese Remainder Theorem): Let  $m_1, m_2, \dots, m_n$  be pairwise relatively prime positive integers greater than 1 and  $a_1, a_2, \dots, a_n$  arbitrary integers. Then, the system

$$x \equiv a_1 \pmod{m_1}$$

$$x \equiv a_2 \pmod{m_2}$$

...

$$x \equiv a_n \pmod{m_n}$$

has a **unique solution** modulo  $m = m_1 m_2 \dots m_n$ .

(That is, there is a solution  $x$  with  $0 \leq x < m$ , and all other solutions are congruent modulo  $m$  to this solution.)



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# The Chinese Remainder Theorem: Example

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

- 1 Let  $m = 3 \cdot 5 \cdot 7 = 105$ ,  $M_1 = m/3 = 35$ ,  $M_2 = m/5 = 21$ , and  $M_3 = m/7 = 15$ .
- 2 Compute  $y_k$ , i.e., the inverse of  $M_k$  modulo  $m_k$ :
  - ▶  $35 \cdot 2 \equiv 1 \pmod{3}$   $y_1 = 2$
  - ▶  $21 \equiv 1 \pmod{5}$   $y_2 = 1$
  - ▶  $15 \equiv 1 \pmod{7}$   $y_3 = 1$
- 3 Compute a solution  $x = a_1M_1y_1 + \dots + a_nM_ny_n$ :  
 $x = 2 \cdot 35 \cdot 2 + 3 \cdot 21 \cdot 1 + 2 \cdot 15 \cdot 1 \equiv 233 \equiv 23 \pmod{105}$
- 4 The solutions are all integers  $x$  that satisfy  $x \equiv 23 \pmod{105}$ .



## Back Substitution

We may also solve systems of linear congruences with pairwise relatively prime moduli  $m_1, m_2, \dots, m_n$  by back substitution.

**Example:**

$$(1) \quad x \equiv 1 \pmod{5}$$

$$(2) \quad x \equiv 2 \pmod{6}$$

$$(3) \quad x \equiv 3 \pmod{7}$$

According to (1),  $x = 5t + 1$ , where  $t$  is an integer.

Substituting this expression into (2), we have  $5t + 1 \equiv 2 \pmod{6}$ , which means that  $t \equiv 5 \pmod{6}$ . Thus,  $t = 6u + 5$ , where  $u$  is an integer.

Substituting  $x = 5t + 1$  and  $t = 6u + 5$  into (3), we have  $30u + 26 \equiv 3 \pmod{7}$ , which implies that  $u \equiv 6 \pmod{7}$ . Thus,  $u = 7v + 6$ , where  $v$  is an integer.

Thus, we must have  $x = 210v + 206$ . Translating this back into a congruence,

$$x \equiv 206 \pmod{210}.$$



# Fermat's Little Theorem

**FERMAT'S LITTLE THEOREM** If  $p$  is prime and  $a$  is an integer not divisible by  $p$ , then

$$a^{p-1} \equiv 1 \pmod{p}.$$

Furthermore, for every integer  $a$  we have

$$a^p \equiv a \pmod{p}.$$



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# RAS Cryptosystem

Pick two large primes  $p$  and  $q$ . Let  $n = pq$ . **Encryption key**  $(n, e)$  and **decryption key**  $(n, d)$  are selected such that

$$(1) \gcd(e, (p-1)(q-1)) = 1$$

$$(2) ed \equiv 1 \pmod{(p-1)(q-1)}$$

**RSA encryption:**  $C = M^e \pmod n$ ;

**RSA decryption:**  $M = C^d \pmod n$ .



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# The Principle of Mathematical Induction

**Well-Ordering Property:** Every nonempty set of nonnegative integers has a least element.

## Principle. (Weak Principle of Mathematical Induction)

(a) **Basic Step:** the statement  $P(b)$  is true

(b) **Inductive Step:** the statement  $P(n - 1) \rightarrow P(n)$  is true for all  $n > b$

Thus,  $P(n)$  is true for all integers  $n \geq b$ .

## Principle (Strong Principle of Mathematical Induction):

(a) **Basic Step:** the statement  $P(b)$  is true

(b) **Inductive Step:** for all  $n > b$ , the statement

$$P(b) \wedge P(b + 1) \wedge \dots \wedge P(n - 1) \rightarrow P(n) \text{ is true.}$$

Then,  $P(n)$  is true for all integers  $n \geq b$ .



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# Recurrence

To specify a function on the basis of a recurrence:

- **Basis step (initial condition)**: Specify the value of the function at zero.
- **Recursive step**: Give a rule for finding its value at an integer from its values at smaller integers.

**Find a closed-form solution?** “Top-down” and “bottom-up”

$$\begin{aligned}T(n) &= rT(n-1) + a \\&= r(rT(n-2) + a) + a \\&= r^2T(n-2) + ra + a \\&= r^2(rT(n-3) + a) + ra + a \\&= r^3T(n-3) + r^2a + ra + a \\&= r^3(rT(n-4) + a) + r^2a + ra + a \\&= r^4T(n-4) + r^3a + r^2a + ra + a.\end{aligned}$$

$$\begin{aligned}T(0) &= b \\T(1) &= rT(0) + a = rb + a \\T(2) &= rT(1) + a = r(rb + a) + a = r^2b + ra + a \\T(3) &= rT(2) + a = r^3b + r^2a + ra + a\end{aligned}$$

Mathematical induction.



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# Counting

**Product Rule:** If a count of elements can be broken down into a **sequence of dependent counts** where the first count yields  $n_1$  elements, the second  $n_2$  elements, and  $k$ -th count  $n_k$  elements, then the total number of elements is

$$n = n_1 \times n_2 \times \dots \times n_k$$

## Sum Rule:

- A task can be done either in one of  $n_1$  ways or in one of  $n_2$  ways
- None of the set of  $n_1$  ways is the same as any of the set of  $n_2$  ways

## The Subtraction Rule:

- A task can be done in **either  $n_1$  ways or  $n_2$  ways**
- **Principle of inclusion–exclusion:**

$$|A_1 \cup A_2| = |A_1| + |A_2| - |A_1 \cap A_2|.$$



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# Pigeonhole Principle

Assume that there are a set of objects and a set of bins to store them.

**The Pigeonhole Principle:** If  $k$  is a positive integer and  $k + 1$  or more objects are placed into  $k$  boxes, then there is at **least one box containing two or more** of the objects.

If  $N$  objects are placed into  $k$  bins, then there is at least one bin containing **at least  $\lceil N/k \rceil$  objects**.



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# Permutations and Combinations

**Theorem:** If  $n$  is a positive integer and  $r$  is an integer with  $1 \leq r \leq n$ , then there are

$$P(n, r) = n(n-1)(n-2) \cdots (n-r+1)$$

$r$ -permutations of a set with  $n$  distinct elements.

**Theorem:** For integers  $n$  and  $r$  with  $0 \leq r \leq n$ , the number of  $r$ -element subsets of an  $n$ -element set is

$$\binom{n}{r} = C(n, r) = \frac{P(n, r)}{r!} = \frac{n!}{r!(n-r)!}$$



# Combinatorial Proof

**Theorem:** Let  $n$  and  $r$  be nonnegative integers with  $r \leq n$ . Then  $C(n, r) = C(n, n - r)$ .

**Definition:** A **combinatorial proof** of an identity is

- a proof that uses counting arguments to prove that **both sides** of the identity **count the same objects** but in different ways
- **or** a proof that is based on showing that there is a **bijection between the sets of objects** counted by the two sides of the identity.

These two types of proofs are called **double counting proofs** and **bijective proofs**, respectively.



# The Binomial Theorem

Let  $x$  and  $y$  be variables, and let  $n$  be a nonnegative integer:

$$(x + y)^n = \sum_{j=0}^n \binom{n}{j} x^{n-j} y^j = \binom{n}{0} x^n + \binom{n}{1} x^{n-1} y + \cdots + \binom{n}{n-1} x y^{n-1} + \binom{n}{n} y^n.$$

**Corollary:** Let  $n$  be a nonnegative integer,

$$\sum_{k=0}^n \binom{n}{k} = 2^n.$$

**Theorem:** Let  $n$  and  $k$  be positive integers with  $n \geq k$ . Then,

$$\binom{n+1}{k} = \binom{n}{k-1} + \binom{n}{k}.$$



# Labelling and Trinomial Coefficients

If we have  $k_1$  labels of one kind (e.g., red),  $k_2$  labels of a second kind (e.g., blue), and  $k_3 = n - k_1 - k_2$  labels of a third kind (e.g., green).

How many different ways to label  $n$  distinct objects?

$$\begin{aligned}\binom{n}{k_1} \binom{n-k_1}{k_2} &= \frac{n!}{k_1!(n-k_1)!} \frac{(n-k_1)!}{(k_2)!(n-k_1-k_2)!} \\ &= \frac{n!}{k_1!k_2!(n-k_1-k_2)!} = \frac{n!}{k_1!k_2!k_3!}\end{aligned}$$

This is called a **trinomial coefficient** and denote it as

$$\binom{n}{k_1 \quad k_2 \quad k_3} = \frac{n!}{k_1!k_2!k_3!},$$

where  $k_1 + k_2 + k_3 = n$ .



# Solving Linear Homogeneous Recurrence Relations

**Definition:** A **linear homogeneous relation** of degree  $k$  with constant coefficients is a recurrence relation of the form

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k},$$

where  $c_1, c_2, \dots, c_k$  are real numbers, and  $c_k \neq 0$ .

By induction, such a recurrence relation is **uniquely** determined by this recurrence relation and  **$k$  initial conditions**  $a_0, a_1, \dots, a_{k-1}$ .





# Solving Linear Homogeneous Recurrence Relations

The characteristic equation (CE) is:

$$r^k - \sum_{i=1}^k c_i r^{k-i} = 0.$$

**Theorem:** Suppose that there are  $t$  roots  $r_1, \dots, r_t$  with multiplicities  $m_1, \dots, m_t$ . Then,

$$\begin{aligned} a_n = & (\alpha_{1,0} + \alpha_{1,1}n + \dots + \alpha_{1,m_1-1}n^{m_1-1})r_1^n \\ & + (\alpha_{2,0} + \alpha_{2,1}n + \dots + \alpha_{2,m_2-1}n^{m_2-1})r_2^n \\ & + \dots + (\alpha_{t,0} + \alpha_{t,1}n + \dots + \alpha_{t,m_t-1}n^{m_t-1})r_t^n \end{aligned}$$

- Solving the roots with CE
- Solving the  $\alpha_i$  for all  $i$  using initial conditions



# Linear Nonhomogeneous Recurrence Relations

**Definition:** A linear nonhomogeneous relation with constant coefficients may contain some terms  $F(n)$  that depend only on  $n$

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n).$$

The recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$  is called the associated homogeneous recurrence relation.

**Theorem:** If  $\{a_n^{(p)}\}$  is any particular solution to the linear nonhomogeneous relation with constant coefficients,

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n),$$

Then all its solutions are of the form

$$a_n = a_n^{(p)} + a_n^{(h)},$$

where  $\{a_n^{(h)}\}$  is any solution to the associated homogeneous recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$ .



# Linear Nonhomogeneous Recurrence Relations

Suppose that  $\{a_n\}$  satisfies the linear nonhomogeneous recurrence relation

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k} + F(n),$$

where  $c_1, c_2, \dots, c_k$  are real numbers, and

$$F(n) = (b_t n^t + b_{t-1} n^{t-1} + \cdots + b_1 n + b_0) s^n,$$

where  $b_0, b_1, \dots, b_t$  and  $s$  are real numbers. When  $s$  is not a root of the characteristic equation of the associated linear homogeneous recurrence relation, there is a particular solution of the form

$$(p_t n^t + p_{t-1} n^{t-1} + \cdots + p_1 n + p_0) s^n.$$

When  $s$  is a root of this characteristic equation and its multiplicity is  $m$ , there is a particular solution of the form

$$n^m (p_t n^t + p_{t-1} n^{t-1} + \cdots + p_1 n + p_0) s^n.$$



# Linear Nonhomogeneous Recurrence Relations

Find all solutions of the recurrence relation  $a_n = 5a_{n-1} - 6a_{n-2} + 7^n$ .

**Solution:**

- $a_n^{(h)} = \alpha_1 \cdot 3^n + \alpha_2 \cdot 2^n$
- Let  $a_n^{(p)} = C \cdot 7^n$ :

$$C \cdot 7^n = 5C \cdot 7^{n-1} - 6C \cdot 7^{n-2} + 7^n.$$

Thus,  $C = 49/20$ , and  $a_n^{(p)} = (49/20)7^n$ .

- Solve  $\alpha_i$  in  $a_n = \alpha_1 \cdot 3^n + \alpha_2 \cdot 2^n + (49/20)7^n$  using initial conditions.



# Generating Function

The **generating function** for the sequence  $a_0, a_1, \dots, a_k, \dots$  of **real numbers** is the infinite series

$$G(x) = a_0 + a_1x + \dots + a_kx^k + \dots = \sum_{k=0}^{\infty} a_kx^k.$$

## Example:

- The sequence  $\{a_k\}$  with  $a_k = 3$

$$\sum_{k=0}^{\infty} 3x^k$$

- The sequence  $\{a_k\}$  with  $a_k = 2^k$

$$\sum_{k=0}^{\infty} 2^k x^k$$



## Generating Function: Finite Series

A finite sequence  $a_0, a_1, \dots, a_n$  can be easily extended by setting  $a_{n+1} = a_{n+2} = \dots = 0$ .

The generating function  $G(x)$  of this infinite sequence  $\{a_n\}$  is a polynomial of degree  $n$ , i.e.,

$$G(x) = a_0 + a_1x + \dots + a_nx^n.$$

**Example:** What is the generating function for the sequence  $a_0, a_1, \dots, a_m$ , with  $a_k = C(m, k)$ ?

$$G(x) = C(m, 0) + C(m, 1)x + C(m, 2)x^2 + \dots + C(m, m)x^m.$$

Based on binomial theorem,  $G(x) = (1 + x)^m$ .

$$(x + y)^n = \sum_{j=0}^n \binom{n}{j} x^{n-j} y^j = \binom{n}{0} x^n + \binom{n}{1} x^{n-1} y + \dots + \binom{n}{n-1} x y^{n-1} + \binom{n}{n} y^n.$$



## Operations of Generating Functions

**Theorem:** Let  $f(x) = \sum_{k=0}^{\infty} a_k x^k$ , and  $g(x) = \sum_{k=0}^{\infty} b_k x^k$ . Then,

$$f(x) + g(x) = \sum_{k=0}^{\infty} (a_k + b_k) x^k$$

$$f(x)g(x) = \sum_{k=0}^{\infty} \left( \sum_{j=0}^k a_j b_{k-j} \right) x^k$$

**Example 2:** To obtain the corresponding sequence of  $G(x) = 1/(1 - ax)^2$  for  $|ax| < 1$ :

Consider  $f(x) = 1/(1 - ax)$  and  $g(x) = 1/(1 - ax)$ . Since the sequence of  $f(x)$  and  $g(x)$  corresponds to  $1, a, a^2, \dots$ , we have

$$G(x) = f(x)g(x) = \sum_{k=0}^{\infty} (k + 1) a^k x^k.$$



# Generating Functions

- For  $|x| < 1$ , function  $G(x) = 1/(1 - x)$  is the generating function of the sequence  $1, 1, 1, 1, \dots$ ,

$$1/(1 - x) = 1 + x + x^2 + \dots$$

- For  $|ax| < 1$ , function  $G(x) = 1/(1 - ax)$  is the generating function of the sequence  $1, a, a^2, a^3, \dots$ ,

$$1/(1 - ax) = 1 + ax + a^2x^2 + \dots$$

- For  $|x| < 1$ ,  $G(x) = 1/(1 - x)^2$  is the generating function of the sequence  $1, 2, 3, 4, 5, \dots$

$$1/(1 - x)^2 = 1 + 2x + 3x^2 + \dots$$





## Example 1

Solve the recurrence relation  $a_k = 3a_{k-1}$  for  $k = 1, 2, 3, \dots$  and initial condition  $a_0 = 2$ .

Let  $G(x)$  be the generating function for the sequence  $\{a_k\}$ , that is,  $G(x) = \sum_{k=0}^{\infty} a_k x^k$ . We aim to first derive the formulation of  $G(x)$ .

$$\begin{aligned}G(x) - 3xG(x) &= \sum_{k=0}^{\infty} a_k x^k - 3 \sum_{k=1}^{\infty} a_{k-1} x^k \\&= a_0 + \sum_{k=1}^{\infty} (a_k - 3a_{k-1}) x^k \\&= 2,\end{aligned}$$

Thus,  $G(x) - 3xG(x) = (1 - 3x)G(x) = 2$ :

$$G(x) = \frac{2}{(1 - 3x)}.$$



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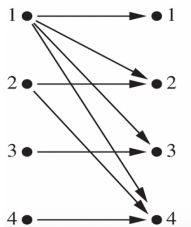
# Cartesian Product

Let  $A = \{a_1, a_2, \dots, a_m\}$  and  $B = \{b_1, b_2, \dots, b_n\}$ , the **Cartesian product**  $A \times B$  is the set of pairs  $\{(a_1, b_1), (a_2, b_2), \dots, (a_1, b_n), \dots, (a_m, b_n)\}$ .

Let  $A$  and  $B$  be two sets. A **binary relation** from  $A$  to  $B$  is a subset of a Cartesian product  $A \times B$ .

A **relation on the set  $A$**  is a relation from  $A$  to **itself**.

We use the notation  $aRb$  to denote  $(a, b) \in R$ , and  $a \not R b$  to denote  $(a, b) \notin R$ .



$R$	1	2	3	4
1	×	×	×	×
2		×		×
3			×	
4				×



# Summary on Properties of Relations

- **Reflexive Relation:** A relation  $R$  on a set  $A$  is called reflexive if  $(a, a) \in R$  for every element  $a \in A$ .
- **Irreflexive Relation:** A relation  $R$  on a set  $A$  is called irreflexive if  $(a, a) \notin R$  for every element  $a \in A$ .
- **Symmetric Relation:** A relation  $R$  on a set  $A$  is called symmetric if  $(b, a) \in R$  whenever  $(a, b) \in R$  for all  $a, b \in A$ .
- **Antisymmetric Relation:** A relation  $R$  on a set  $A$  is called antisymmetric if  $(b, a) \in R$  and  $(a, b) \in R$  implies  $a = b$  for all  $a, b \in A$ .
- **Transitive Relation:** A relation  $R$  on a set  $A$  is called transitive if  $(a, b) \in R$  and  $(b, c) \in R$  implies  $(a, c) \in R$  for all  $a, b, c \in A$ .



# Combining Relations

**Definition:** Let  $R$  be a relation from a set  $A$  to a set  $B$  and  $S$  be a relation from  $B$  to  $C$ . The composite of  $R$  and  $S$  is the relation consisting of the ordered pairs  $(a, c)$  where  $a \in A$  and  $c \in C$  and for which there is a  $b \in B$  such that  $(a, b) \in R$  and  $(b, c) \in S$ .

**Example:** Let  $A = \{1, 2, 3\}$ ,  $B = \{0, 1, 2\}$ , and  $C = \{a, b\}$ :

- $R = \{(1, 0), (1, 2), (3, 1), (3, 2)\}$
- $S = \{(0, b), (1, a), (2, b)\}$
- $S \circ R = \{(1, b), (3, a), (3, b)\}$



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