

ECE 120: Introduction to Computing

Sparse Representations, Bit Errors, and Error Detection

Representations Must be Unambiguous

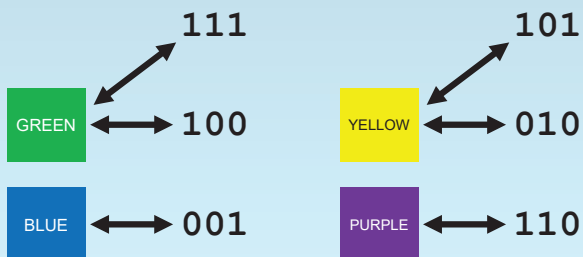
Recall requirements for a representation...

A bit pattern represents **at most one thing**.



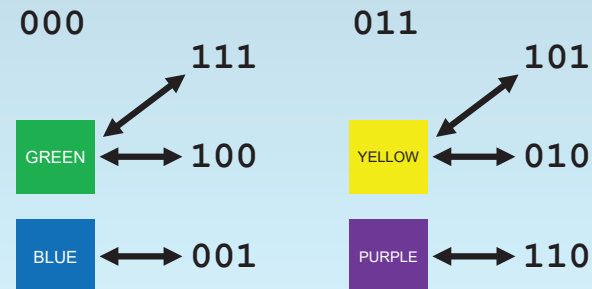
But One Thing Can Be Represented by Several Patterns

For example, **100** and **111**
can both represent **GREEN**.



And Some Patterns May Represent Nothing

Bit patterns **000** and **011** represent no color.



Example: Binary-Coded Decimal (BCD)

BCD maps **decimal digits** into 4-bit patterns.

Each pattern is the binary value of the digit.

Six patterns are unused ($2^4 - 10 = 6$).

0	↔	0000	5	↔	0101	1010 1011 1100 1101 1110 1111
1	↔	0001	6	↔	0110	
2	↔	0010	7	↔	0111	
3	↔	0011	8	↔	1000	
4	↔	0100	9	↔	1001	

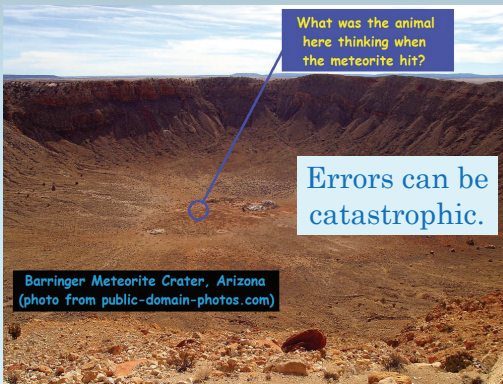
Unused Bit Patterns Can Be Used to Detect Errors

Bit patterns on the right have no meaning.

Seeing one of these patterns means that **an error has occurred**.

0	↔	0000	5	↔	0101	1010 1011 1100 1101 1110 1111
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But What Exactly is an “Error?”



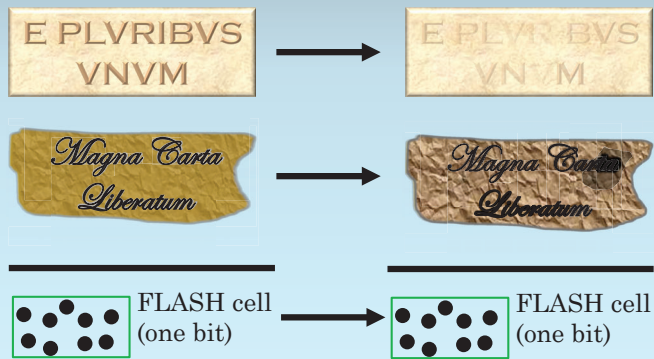
We Need to Choose a Model for Errors

Cannot handle “all” errors.

- Catastrophes are always possible.
- But catastrophes are (hopefully) uncommon!

Let’s focus on frequent types of errors.

Errors are Not a New Phenomenon



Our Class Uses Bit Flips as an Error Model

Previous slide showed **erasures**, in which a character is missing.

our class assumes

- a more difficult type of error
 - one character is replaced by another
- with binary representations
- these are **bit flips**
 - **0 becomes 1, or 1 becomes 0**



Examples from English: Erasures vs. Errors

·hat ·oes t·is ·ay·

What does this say?

Than foen thit xay.

That goes this way.

Assume Small Numbers of Independent Bit Flips

Bit flip error model: 0 becomes 1, 1 becomes 0

Each bit flips

- independently of all others
- with some low probability (call it p)

For N bits

- chance of one error is $Np(1-p)^{N-1}$
- chance of two errors is $\frac{1}{2} N(N-1)p^2(1-p)^{N-2}$

In practice, $Np \ll 1$ (Np is much less than 1) so
chance of two errors \ll chance of one error

Example: 2-out-of-5 Code

A 2-out-of-5 code maps **decimal digits** into 5-bit patterns.

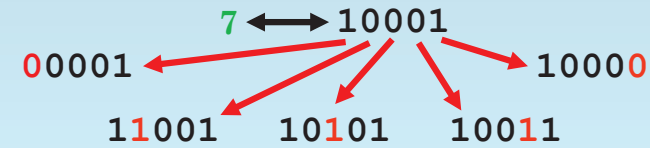
Each pattern has **exactly two 1 bits**.

1	↔	00011	6	↔	01100
2	↔	00101	7	↔	10001
3	↔	00110	8	↔	10010
4	↔	01001	9	↔	10100
5	↔	01010	0	↔	11000

Claim: Can Always Detect a Single Bit Error

Thought experiment:

What happens if a single bit flips in a digit coded using a 2-out-of-5 code?



None of these patterns means anything, so **we can always detect the error!**