University of Illinois at Urbana-Champaign
Dept. of Electrical and Computer Engineering
ECE 120: Introduction to Computing

Bit-Sliced Comparator

## How Do You Compare Unsigned Numbers?

Let's develop a bit-sliced design to compare two unsigned numbers.

Which 8-bit unsigned number is bigger?

01101000
01010111
How did you know?
Did you start on the left or the right?

ECE 120: Introduction to Computing $\quad$ O2016 Steven S. Lumetta. All rights reserved.
slide 2

## Humans Go from Left to Right

Usually, humans start on the left. Why?
As soon as we notice a difference, we're done!

## humans compare this way

$$
\begin{aligned}
& \hline 011 \mathrm{a}_{4} \mathrm{a}_{3} \mathrm{a}_{2} \mathrm{a}_{1} \mathrm{a}_{0} \\
& 010 \mathrm{~b}_{4} \mathrm{~b}_{3} \mathrm{~b}_{2} \mathrm{~b}_{1} \mathrm{~b}_{0}
\end{aligned}
$$

Bit-sliced hardware cannot stop in the middle.
The information flows from one end to the other.
Output wires produce the answer (in bits).

## Our Design Compares from Right to Left

Our comparator design will start on the right.
humans compare this way
$\overrightarrow{a_{7} a_{6} a_{5} a_{4} a_{3} a_{2} a_{1} a_{0}}$
$\mathrm{b}_{7} \mathrm{~b}_{6} \mathrm{~b}_{5} \mathrm{~b}_{4} \mathrm{~b}_{3} \mathrm{~b}_{2} \mathrm{~b}_{1} \mathrm{~b}_{0}$
our design will compare this way
From least significant to most significant bit.

## Three Possible Answers for Comparison of A and B

When comparing two numbers, A and B , we have three possible outcomes:

$$
\begin{aligned}
& \mathrm{A}<\mathrm{B} \\
& \mathrm{~A}=\mathrm{B} \\
& \mathrm{~A}>\mathrm{B}
\end{aligned}
$$

To decide the answer for $\mathrm{N}+1$ bits, we need: - the answer for $\mathbf{N}$ (less significant) bits,

- one bit of A, and
${ }^{\circ}$ one bit of $\mathbf{B}$.


## An Abstract Model of the Comparator Bit Slice

A question for you:
How many bits must pass between slices?
Two!
This figure shows an abstract model of our bit slice.


ECE 120: Introduction to Computing $\quad \bigcirc 2016$ Steven S. Lumetta. All rights reserved.
slide 6

## We Need a Representation for Answers

Another question for you:
How do we represent the three possible answers?
Any way we want!
Our choice of representation will affect the amount of logic we need.
Here's a good one...

| $\mathbf{C}_{1}$ | $\mathbf{C}_{\mathbf{0}}$ | meaning |
| :---: | :---: | :---: |
| 0 | 0 | A $=\mathrm{B}$ |
| 0 | 1 | A $<\mathrm{B}$ |
| 1 | 0 | A $>\mathrm{B}$ |
| 1 | 1 | not used |

## A Single Bit Requires Two Minterms on A, B

Let's start by solving a single bit.
In this case, there are no less significant bits.
So we consider
only A and B.
Fill in the meanings,
then the bits.

| A | $\mathbf{B}$ | $\mathrm{Z}_{1}$ | $\mathrm{Z}_{0}$ | meaning |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{0}$ | $\mathbf{0}$ | 0 | 0 | $\mathrm{~A}=\mathrm{B}$ |

Note that $Z_{1}$ and $Z_{0}$ are minterms.
$\begin{array}{lllll}1 & 0 & 1 & 0 & \mathrm{~A}>\mathrm{B} \\ 1 & 1 & 0 & 0 & \mathrm{~A}=\mathrm{B}\end{array}$

## Comparing Two Bits is Fairly Easy

An implementation for a single bit appears below.
This structure forms the core of our bit slice, since it compares one bit of A with one bit of B.


ECE 120: Introduction to Computing © 2016 Steven S. Lumetta. All rights reserved.

## When A and B are Equal, Pass Along the Answer

Now for the full problem.
We'll start with the case of $\mathbf{A}=0$ and $\mathbf{B}=0$.

| A B | $\mathbf{C}_{1}$ | $\mathbf{C}_{0}$ | meaning | $\mathrm{Z}_{1}$ | $\mathrm{Z}_{0}$ | meaning |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | $\mathrm{~A}=\mathrm{B}$ | 0 | 0 | $\mathrm{~A}=\mathrm{B}$ |
| 0 | 0 | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ |
| 0 | 0 | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ |
| 0 | 0 | 1 | 1 | ??? | $\mathbf{x}$ | $\mathbf{x}$ | don't care |

ECE 120: Introduction to Computing $\quad \bigcirc 2016$ Steven S. Lumetta. All rights reserved.
slide 10

## When A and B are Equal, Pass Along the Answer

Is there any difference when $\mathrm{A}=1$ and $\mathrm{B}=1$ ?
No, outputs are the same as the last case.

| A | B | $\mathrm{C}_{1}$ | $\mathrm{C}_{0}$ | meaning | $\mathrm{Z}_{1}$ | $\mathrm{Z}_{0}$ | meaning |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 0 | 0 | $\mathrm{~A}=\mathrm{B}$ | 0 | 0 | $\mathrm{~A}=\mathrm{B}$ |
| 1 | 1 | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ |
| 1 | 1 | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ |
| 1 | 1 | 1 | 1 | ??? | $\mathbf{x}$ | $\mathbf{x}$ | don't care |

## When A and B Differ, Override the Previous Answer

What about case of $\mathrm{A}=0$ and $\mathrm{B}=1$ ?
Always output $A<B$ (for valid inputs).

| A | B | $\mathrm{C}_{1}$ | $\mathrm{C}_{0}$ | meaning | $\mathrm{Z}_{1}$ | $\mathrm{Z}_{0}$ | meaning |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 0 | $\mathrm{~A}=\mathrm{B}$ | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ |
| 0 | 1 | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ |
| 0 | 1 | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ |
| 0 | 1 | 1 | 1 | ??? | $\mathbf{x}$ | $\mathbf{x}$ | don't care |

## When A and B Differ, Override the Previous Answer

And the case of $\mathrm{A}=1$ and $\mathrm{B}=0$ ?
Always output A > B (for valid inputs).

| A | B | $\mathbf{C}_{1}$ | $\mathbf{C}_{0}$ | Meaning | $\mathrm{Z}_{1}$ | $\mathrm{Z}_{0}$ | meaning |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 0 | 0 | 0 | $\mathrm{~A}=\mathrm{B}$ | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ |
| 1 | 0 | 0 | 1 | $\mathrm{~A}<\mathrm{B}$ | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ |
| 1 | 0 | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ | 1 | 0 | $\mathrm{~A}>\mathrm{B}$ |
| 1 | 0 | 1 | 1 | $? ? ?$ | $\mathbf{x}$ | $\mathbf{x}$ | don't care |

ECE 120: Introduction to Computing

## $\mathrm{Z}_{1}$ is a Majority Function

Let's use a K-map to solve $\mathbf{Z}_{1}$.
What are the loops?
AB'
$\mathrm{AC}_{1}$
$\mathrm{B}^{\prime} \mathrm{C}_{1}$
$\stackrel{\mathrm{So}}{\mathrm{Z}_{1}}=\mathrm{AB}^{\prime}+\mathrm{AC}_{1}+\mathrm{B}^{\prime} \mathrm{C}_{1}$

| $\mathrm{Z}_{1}$ | AB |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | 00 | 01 | 11 | 10 |
| 00 | 0 | 0 | 0 | 1 |
| 01 | 0 | 0 | 0 | 1 |
| $\mathrm{C}_{1} \mathrm{C}_{0}$ | $x$ | x | x | x |
| 10 | 1 | 0 | 1 | 1 |

ECE 120: Introduction to Computing
© 2016 Steven S. Lumetta. All rights reserved.
slide 14

## Full Implementation as SOP Expressions



