University of Illinois at Urbana-Champaign
Dept. of Electrical and Computer Engineering

## ECE 120: Introduction to Computing

## 2's Complement Overflow and <br> Boolean Logic

## Example: Addition of Unsigned Bit Patterns

A question for you:
What is the overflow condition for addition of two N-bit 2's complement bit patterns?
(That is, when is the sum incorrect?)
Remember that addition works exactly the same way as with N-bit unsigned bit patterns, so we can do some base 2 addition to find the answer.

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## Adding Two Non-Negative Patterns Can Overflow

Let's start with our first example from before:

$$
\begin{aligned}
& 11 \\
& 01110(14) \\
& +00100(4) \\
& \hline 10010(-14)
\end{aligned}
$$

Oops! We had no carry out, but the answer is wrong (an overflow occurred).

So overflow is different than for unsigned...

## Carry Out Does Not Indicate 2's Complement Overflow

This example overflowed when the bits were interpreted with an unsigned representation.

```
We have no (2)11
    space for 01110 (14)
        that bit! 01110 (14)
    + 10101 (-11;21 unsigned)
    00011 (3)
```

But here the answer is still correct!
Carry out $\neq$ overflow for 2 's complement.

## Adding Non-Negative to Negative Can Never Overflow

## Claim:

Addition of two N-bit 2's complement bit patterns can not overflow if one pattern is negative (starts with 1 ) and the other pattern is non-negative (starts with 0 ).
Proof: You do it!
And THEN you can read the proof in the notes.

## Long Definition for Overflow of 2's Complement Addition

Add two N-bit 2's complement patterns.
A $a_{N-2} \ldots a_{0}$ (sign bit is $A$ )
$+B b_{N-2} \ldots b_{0}$ (sign bit is B)
$\mathbf{S} \mathbf{S}_{\mathrm{N}-2} \ldots \mathbf{S}_{0}$ (sign bit is S )
Claim: The addition overflows iff one of the following holds:

1. The two addends are non-negative, and the sum is negative.
2. The two addends are negative, and the sum is non-negative.

## Boolean Algebra Gives a More Concise Expression

That's a lot of words!
Boolean algebra gives a more concise form:
OVERFLOW =
[ (NOT A) AND (NOT B) AND S ] OR
[ A AND B AND (NOT S) ]
(Remember: A, B, and S were the sign bits.)
But what do these operators (AND, OR, and NOT) mean?

## Boolean Operators Were Invented in the mid-19 ${ }^{\text {th }}$ Century

Boolean operators were invented (by George Boole) to reason about logical propositions.
They originally operated on true/false values.
We use them with ... that's right, bits!

$$
0=\text { false and } 1=\text { true }
$$

Be careful not to confuse Boolean operators with English words. The meanings are not identical.

## We Use Only a Few Boolean Functions

AND: the ALL function
returns 1 iff ALL inputs are 1 (otherwise 0 )
OR: the ANY function
returns 1 iff ANY input is 1 (otherwise 0 )
NOT: logical complement
(NOT 0 ) is 1 ; (NOT 1 ) is 0
XOR: the ODD function
returns 1 iff an ODD number of inputs
are 1 (otherwise 0)

## A Truth Table Fully Defines a Boolean Function

The drawing to the right is a truth table.
A truth table allows us to - define a Boolean function C - by listing the output value - for all combinations of inputs (here A and B, in base 2 order).

| $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ |
| :--- | :--- | :--- |
| 0 | 0 |  |
| 0 | 1 |  |
| 1 | 0 |  |
| 1 | 1 |  |

Let's write truth tables for our four Boolean functions.

## AND: The ALL Function



## OR: The ANY Function

| And now OR. | A B | A OR B |
| :---: | :---: | :---: |
| OR can also be written in other ways: | 00 | 0 |
|  | 01 | 1 |
| $\begin{array}{l\|l} \circ \mathrm{A}+\mathrm{B} & \begin{array}{l} \text { We usually } \\ \text { use this one. } \end{array} \end{array}$ | $\begin{array}{ll}1 & 1 \\ 1 & 0\end{array}$ | 1 |
| - $\mathbf{A} \vee \mathrm{B}$ (math. disjunction) | 11 | 1 |
| Note rounded in pointed out |  | -A+B |

## NOT: Logical Complement

| And now NOT. |  | A | NOT A |
| :---: | :---: | :---: | :---: |
| NOT can also be written in other ways: |  | 0 | 1 |
|  |  | 1 | 0 |
| $\stackrel{\circ}{\circ}{ }^{\circ}{ }^{\prime} \left\lvert\, \begin{array}{ll}\text { We usually } \\ \text { use these. }\end{array}\right.$ <br> - $\neg$ A (math. complement) |  |  |  |
|  |  |  |  |

## XOR: The ODD Function

| And, finally, XOR. | A | $\mathbf{B}$ | A XOR B |
| :--- | :---: | :---: | :---: |
| XOR is usually written <br> this way: $A \oplus B$ | 0 | 0 | 0 |
|  | 0 | 1 | 1 |
| 1 | 0 | 1 |  |
|  | 1 | 1 | 0 | double line for inputs.

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Use Definitions to Generalize to More than Two Operands

| Generalize to more | $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{A} \oplus \mathbf{B} \oplus \mathbf{C}$ |
| :--- | :---: | :---: | :---: | :---: |
| operands using the <br> definitions given: | 0 | 0 | 0 | 0 |
| - AND: ALL | 0 | 0 | 1 | 1 |
| $\circ$ OR: ANY | 0 | 1 | 0 | 1 |
| $\circ$ XOR: ODD | 0 | 1 | 1 | 0 |
| As an example, fill the | 1 | 0 | 0 | 1 |
| truth table for a | 1 | 0 | 1 | 0 |
| 3-input XOR. | 1 | 1 | 0 | 0 |
|  | 1 | 1 | 1 | 1 |

## Generalize to Sets of Bits by Pairing Bits

We can also generalize to sets of bits.
For example, if we have two N -bit patterns,

$$
\mathrm{A}=\mathrm{a}_{\mathrm{N}-1} \ldots \mathrm{a}_{0} \text { and } \mathrm{B}=\mathrm{b}_{\mathrm{N}-1} \ldots \mathrm{~b}_{0}
$$

we can write

$$
\mathrm{C}=\mathrm{A} \text { AND } \mathrm{B}
$$

To mean that

$$
\text { if } \mathrm{C}=\mathrm{c}_{\mathrm{N}-1} \ldots \mathrm{c}_{0}, \mathrm{c}_{\mathrm{i}}=\mathrm{a}_{\mathrm{i}} \mathrm{~b}_{\mathrm{i}} \text { for } 0 \leq \mathrm{i}<\mathrm{N} \text {. }
$$

## Don't Mix Algebras: Use AND/OR/NOT for Bitwise Logic

If A is a 2 's complement bit pattern, we might also write $-\mathrm{A}=($ NOT A $)+1$
Be careful about mixing

- algebraic notation for Boolean functions - with arithmetic operations.

The " + " in the equation above means base 2 addition (and discarding any carry out), not OR.

