## ECE 120 Final Exam

Fall 2016

Wednesday, December 14, 2016


- Be sure that your exam booklet has 14 pages.
- Write your name, netid and check discussion section on the title page.
- Do not tear the exam booklet apart, except for the last four pages.
- Use backs of pages for scratch work if needed.
- This is a closed book exam. You may not use a calculator.
- You are allowed two handwritten $8.5 \times 11$ " sheets of notes (both sides).
- Absolutely no interaction between students is allowed.
- Clearly indicate any assumptions that you make.
- The questions are not weighted equally. Budget your time accordingly.

Problem 120 points
Problem 216 points $\qquad$
Problem 314 points $\qquad$
Problem 421 points $\qquad$
Problem 514 points $\qquad$
Problem 68 points $\qquad$
Problem $7 \quad 7$ points $\qquad$
$\qquad$

Problem 1 (20 points): Binary Representation and Operations, Hamming codes

1. (2 points) There are 365 days in a year. If we want to uniquely identify each day using 2's complement binary representation, what is the minimum number of bits we should use?

Minimum number of bits: $\qquad$ (decimal number)
2. (4 points) Convert the following 24-bit pattern to hexadecimal:
$11000000111111111110 \mathbf{1 1 1 0}_{2}=x$ $\qquad$ (hexadecimal number)
3. (4 points) Perform the following bitwise logical operations.
a) 0110 NAND $0011=$ $\qquad$
b) $1001 \operatorname{XOR}(\operatorname{NOT}(0101))=$ $\qquad$
4. (4 points) Perform the following operation in four-bit 2's complement representation.
$0101+101=$ $\qquad$
Circle one: Carry out? YES NO

Circle one: Overflow? YES NO
5. (6 points) Someone just sent you the following 7-bit Hamming code:
$\mathrm{X}_{7} \mathrm{X}_{6} \mathrm{X}_{5} \mathrm{X}_{4} \mathrm{X}_{3} \mathrm{X}_{2} \mathrm{X}_{1}=1010111$. Does the message have an error or not?

## Circle one: YES NO

If you think there is an error, write the position where there is an error:

There is an error in position $\qquad$

## Problem 2 (16 points): LC-3 Assembly Programming

Greetings, ECE 120 student.
Your mission, should you choose to accept it, is to write the missing lines of code, so the program can properly print on screen a message to wish you an enjoyable break. Additionally, you must write the missing entries in the symbol table associated with this program.
As always, should you or any of your friends be caught or killed, the ECE 120 instructors will disavow any knowledge of your actions. This page will self-destruct by the end of the semester.

Good luck, ECE 120 student.

1. (11 points) Write the missing lines of code. You must write one instruction per missing line.
.ORIG x6000

|  |  | ; Print "Choose message: |
| :---: | :---: | :---: |
|  |  | ; |
|  | LD R1, OPTION | ; R1 <- M[OPTION] |
|  |  | ; Read from keyboard |
|  | NOT R0, R0 | ; R0 <- R1-R0 |
|  | - ; | ; |
|  | ADD R0, R1, R0 | ; |
|  |  | ; Character typed $=$ R1? |
|  | ; Case: character typed = R1 |  |
| EQUAL | LEA RO, HOLIDAYS ; | ; RO <- HOLIDAYS |
|  | BRnzp PRINTOUT | ; Go to PRINTOUT |
|  | ; Case: character typed $\neq$ R1 |  |
| DIFFERENT | LEA R0, NEWYEAR | ; RO <- NEWYEAR |
| PRINTOUT | PUTS | ; Print selected message |
|  |  | ; |
| PROMPT | . STRINGZ "Choose message: " |  |
| OPTION | .FILL x0031 | ; ASCII '1' |
| HOLIDAYS | .STRINGZ "Happy Holidays!" |  |
| NEWYEAR | .STRINGZ "Happy New Year!" |  |
|  | .END |  |

## Problem 2 (16 points): LC-3 Assembly Programming, continued

2. (5 points) Write the missing entries in the symbol table. Answers in hexadecimal only.

| // Symbol table |  |  |
| :---: | :---: | :---: |
| // Scope level 0: |  |  |
| / / | Symbol Name | Page Address |
| / / |  |  |
| / / | EQUAL |  |
| / / | DIFFERENT | 600A |
| / / | PRINTOUT | 600B |
| / / | PROMPT | 600 D |
| / / | OPTION |  |
| / / | HOLIDAYS | 601 F |
| / / | NEWYEAR |  |

## Problem 3 (14 points): Synchronous Counter

1. (11 points) Using $D$ flip-flops, design a 3-bit counter that counts the prime number sequence $2,3,5,7$, and repeats. The current state of the counter is denoted by $\mathrm{S}_{2} \mathrm{~S}_{1} \mathrm{~S}_{0}$. Fill in the K-maps for $\mathrm{S}_{2}{ }^{+}, \mathrm{S}_{1}{ }^{+}$and $\mathrm{S}_{0}{ }^{+}$using don't cares wherever possible.


Write minimal SOP Boolean expressions for $\mathrm{S}_{2}{ }^{+}, \mathrm{S}_{1}{ }^{+}$, and $\mathrm{S}_{0}{ }^{+}$.

$$
\begin{aligned}
& \mathrm{S}_{2}{ }^{+}= \\
& \mathrm{S}_{1}{ }^{+}= \\
& \mathrm{S}_{0}{ }^{+}= \\
&
\end{aligned}
$$

2. (3 points) Suppose you have already designed a 2-bit binary up-counter that counts in the sequence $0,1,2,3$, and repeats. You could attach output logic so that the 2-bit state of this counter produces a 3 -bit output: the repeating prime number sequence $2,3,5,7$. Write down one advantage of the approach described here compared to the implementation in part 1. Express your answer in 10 words or fewer. (We will not read more than 10 words.)

## Problem 4 (21 points): LC-3 Data Path and Control Unit

1. (12 points) The registers of an LC-3 processor have the values shown below to the right.

Consider the LC-3 instructions shown in the table below. For the execute state of each instruction (state number is provided), fill in the values in the instruction register (IR), at the A input of the ALU, at the B input of the ALU, and on the bus. Write all answers in hexadecimal.

| R4 | $x 4444$ |
| :--- | :---: |
| R5 | $x 5555$ |
| R6 | $x 6666$ |
| R7 | $x 7777$ |
|  |  |


| $c$ | State | IR | A input <br> of ALU |  |  | B input <br> of ALU |  | Bus |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Instruction | number | R1, R5, R5 | 5 | $\times 5345$ | $\times 5555$ | $\times 5555$ |  |  |
| ADD R0, R4, \#8 | 1 |  |  |  |  |  |  |  |
| NOT R2, R7 | 9 |  |  |  |  |  |  |  |

2. (9 points) Suppose the LC-3 designers redefine the BR instruction. The 16-bit format stays the same, but the new RTL (after fetch and decode phases) is:

## $B E N: \quad P C \leftarrow M[P C+S E X T(P C o f f s e t 9)]$

In other words, if BEN=1 then PC changes. Complete the LC-3 FSM diagram below. Fill in the four states for BR with RTL, and draw state transitions with labels (if appropriate). Do NOT number the states.


## Problem 5 (14 points): FSM Analysis

The FSM on the left below performs a serial calculation on an input A. Four bits are provided through A each cycle. In the first cycle, the F input ("first bits") is set to 1 . In all subsequent cycles, $\mathrm{F}=0$. After N cycles, the value S provides the answer as an unsigned number.

The size of the FSM depends on the parameter k, which must be at least 3 . Notice that the FSM makes use of a register to hold the state ( $S$ is just the stored register value), a set of k 2-to-1 muxes controlled by F, and a k-bit adder. The mystery box (implementation shown on the right below) transforms $A$ into a 3-bit value $B$, which is then treated as an unsigned number and zero-extended (padded with leading 0 s ) to k bits.


The questions you need to answer are in the following page.
Tear the last page and use it as scratch paper.

## Problem 5 (14 points): FSM Analysis, continued

Answer the questions below based on the FSM design and description on the previous page. In order to help you solving these questions, we strongly suggest that you fill in the truth table for the mystery box. To do that, feel free to tear apart the last page of the exam and use it as scratch paper, because we will NOT grade the truth table.

Circle EXACTLY ONE ANSWER for each question.

1. (3 points) What is the smallest possible value represented by the unsigned bit pattern $B$, given the implementation of the mystery box?
a) -4
b) 4
c) 1
d) -3
e) 0
2. ( 3 points) What is the largest possible value represented by the unsigned bit pattern $B$, given the implementation of the mystery box?
a) 7
b) 0
c) 3
d) 4
e) -4
3. (4 points) The V output from the adder signifies overflow in the stored value. In terms of $k$, what is the minimum number of cycles (including the F=1 cycle) for which the FSM can execute before $\mathrm{V}=1$ ?
a) 1
b) $2^{k-2}-1$
c) $2^{k-1}-1$
d) $1-2^{k}$
e) ceil $\left(2^{k} / 7\right)-1$
4. (4 points) What is the meaning of the output $S$ ?
a) $S$ is the number of cycles in which input $A$ has an odd number of 1 bits.
b) $S$ is the number of 1 bits passed in through $A$.
c) $S$ is the sum of 2's complement values passed in through $A$.
d) $S$ is the number of 0 bits passed in through $A$.
e) None of the above.

## Problem 6 (8 points): LC-3 Instructions and Assembler

1. ( 5 points) Decode each of the following LC-3 instructions, writing the RTL in the box beside the instruction. For full credit, your RTL must include specific values for each operand (for example, "R4" rather than "DR"), and must be sign-extended when appropriate. Do not perform calculations such as addition of the PC value.

You may write any immediate values either as hexadecimal (prefix them with " $x$ ") or as decimal (prefix them with "\#").

Hint: Draw lines between bits to separate the instructions into appropriate fields.

| Instruction bits |  |  |  | RTL Meaning |
| :---: | :---: | :---: | :---: | :---: |
| 0001 | 1110 | 1011 | 0010 | R7 $\leftarrow$ R2 - \#14, setcc |
| 1100 | 0001 | 0100 | 0000 |  |
| 1011 | 0010 | 0101 | 0011 |  |
| 0110 | 0010 | 1000 | 0011 |  |

2. (3 points) The LC-3 assembler finds a single error in the following code. State the nature of the error and in which pass the assembler identifies the error (first or second).
```
    .ORIG x3000
    LEA R1,STRING
PRINT LDR R0,R1,#0
    BRz DONE
    TRAP x21 ; OUT
    ADD R1,R1,#1
    BRnzp PRINT
DONE LEA R1,STRING
AGAIN LDR R0,R1,#0
    BRz DONE
    TRAP x21 ; OUT
    ADD R1,R1,#1
    BRnzp AGAIN
    DONE HALT
STRING .STRINGZ "This is my string."
DATA .FILL xFFFF
.END
```

Circle one: PASS $1 \quad$ PASS 2
Nature of error:
Express your answer in 10 words or fewer. (We will not read more than 10 words.)

## Problem 7 (7 points): LC-3 Assembly Language Interpretation

All questions for this problem pertain to the following code.

```
    .ORIG x3000
    LDI R1,MAGIC
    AND R3,R3,#0
OUTER AND R2,R2,#0 ; outer loop starts here
    AND R0,R0,#0
INNER ADD R0,R0,R0 ; inner loop starts here
    ADD R1,R1,#0 ; the inner loop left shifts bits R1[15:12]
    BRzp ZEROBIT ; out of R1 and into R0[3:0] to form
    ADD RO,RO,#1 ; a single hex digit
ZEROBIT ADD R1,R1,R1
    ADD R2,R2,#1
    ADD R4,R2,#-4
    BRn INNER ; end of inner loop
    ADD R4,R0,#-10 ; start of 'curious code'
    BRzp FORWARD
    LD R2,DIGITO
    ADD R0,R0,R2
    BRnzp LABEL
FORWARD LD R2,LETTERA
    ADD R0,R4,R2
LABEL OUT ; end of 'curious code'
    ADD R3,R3,#1
    ADD R4,R3,#-4
    BRn OUTER ; end of outer loop
    LD RO,NEWLN
    OUT
    HALT
MAGIC .FILL x4000
DIGITO .FILL x30 ; ASCII digit 0 ('0')
LETTERA .FILL x41 ; ASCII letter A ('A')
NEWLN .FILL xOA ; ASCII newline character ('\n')
    .END
```

1. (1 point) How many times does the body of the outer loop execute?
2. (1 point) How many times does the body of the inner loop execute (for each outer loop iteration)?
3. (3 points) What does the 'curious code' marked in the comments do? Express your answer in 10 words or fewer. (We will not read more than 10 words.)
4. (2 points) Explain how to make the program print "ECEB" followed by a newline character to the LC-3 display. Express your answer in 10 words or fewer. (We will not read more than 10 words.)

LC-3 TRAP Service Routines


## LC-3 Control Word Fields



## LC-3 Microsequencer Control

## Signal

Description
$\operatorname{IRD}\left\{\begin{array}{l}=1, \mathrm{CAR} \leftarrow 00 \mid l o p c o d e \\ =0\end{array}\right.$ (opcode $=\operatorname{IR}[15: 12]$ ), only during decode , CAR $\leftarrow J$ (plus 1,2,4,0,16 depending on COND bits)

J 6-bit next value for CAR (plus modifications depending on COND bits)


LC-3 Instructions



Problem 5's help page (use as scratch copy, we will NOT grade it)


| $A_{3}$ | $A_{2}$ | $A_{1}$ | $A_{0}$ | $\mathrm{~B}_{2}$ | $\mathrm{~B}_{1}$ | $\mathrm{~B}_{0}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |  |  |  |
| 0 | 0 | 0 | 1 |  |  |  |
| 0 | 0 | 1 | 0 |  |  |  |
| 0 | 0 | 1 | 1 |  |  |  |
| 0 | 1 | 0 | 0 |  |  |  |
| 0 | 1 | 0 | 1 |  |  |  |
| 0 | 1 | 1 | 0 |  |  |  |
| 0 | 1 | 1 | 1 |  |  |  |
| 1 | 0 | 0 | 0 |  |  |  |
| 1 | 0 | 0 | 1 |  |  |  |
| 1 | 0 | 1 | 0 |  |  |  |
| 1 | 0 | 1 | 1 |  |  |  |
| 1 | 1 | 0 | 0 |  |  |  |
| 1 | 1 | 0 | 1 |  |  |  |
| 1 | 1 | 1 | 0 |  |  |  |
| 1 | 1 | 1 | 1 |  |  |  |

## REPLICATED FROM PROBLEM STATEMENT FOR YOUR CONVENIENCE:

The FSM on the left performs a serial calculation on an input A. Four bits are provided through A each cycle. In the first cycle, the F input ("first bits") is set to 1 . In all subsequent cycles, $\mathrm{F}=0$. After N cycles, the value $S$ provides the answer as an unsigned number.

The size of the FSM depends on the parameter k , which must be at least 3. Notice that the FSM makes use of a register to hold the state ( S is just the stored register value), a set of k 2-to-1 muxes controlled by F, and a k-bit adder. The mystery box (implementation shown above on the left) transforms A into a 3-bit value B, which is then treated as an unsigned number and zero-extended (padded with leading 0 s ) to k bits.

